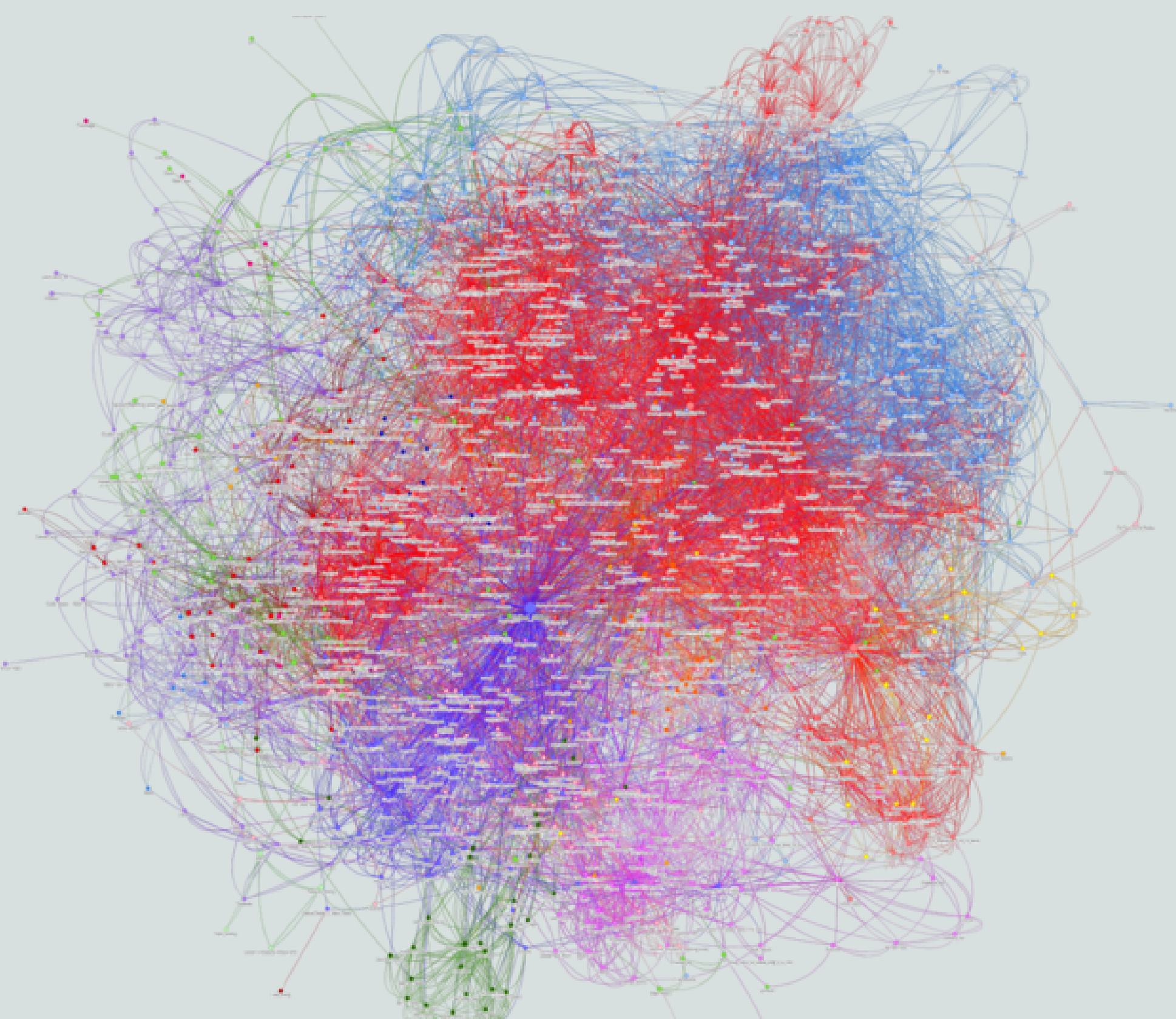


CLIQUE PERCOLATION METHOD: MEMORY EFFICIENT ALMOST EXACT COMMUNITIES

Automatic detection of relevant groups of nodes in large real-world graphs, i.e. community detection, has applications in many fields and has received a lot of attention in the last twenty years. The most popular method designed to find overlapping communities (where a node can belong to several communities) is perhaps the Clique Percolation Method (CPM).



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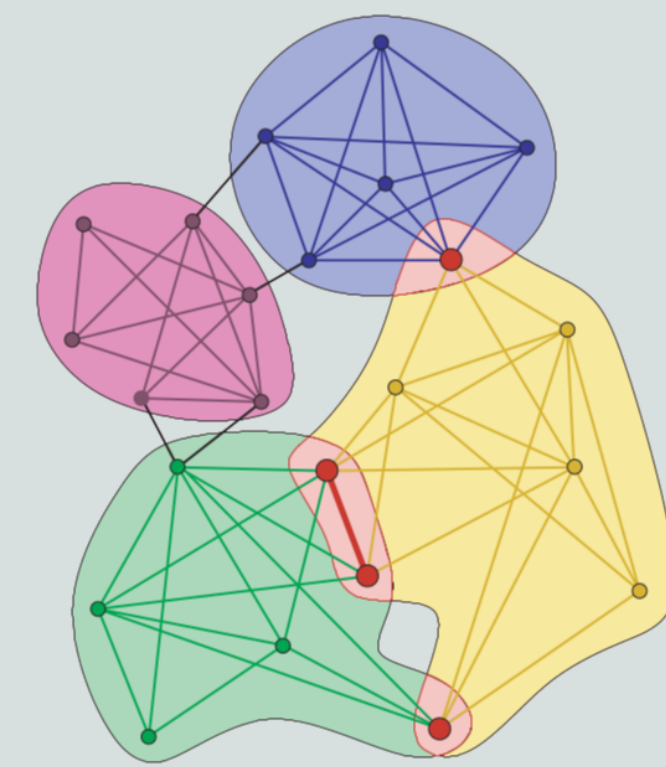


CONTEXT

COMMUNITIES IN A GRAPH

Set of nodes:

- Densely connected *inside*;
- Sparsely connected *outside*.



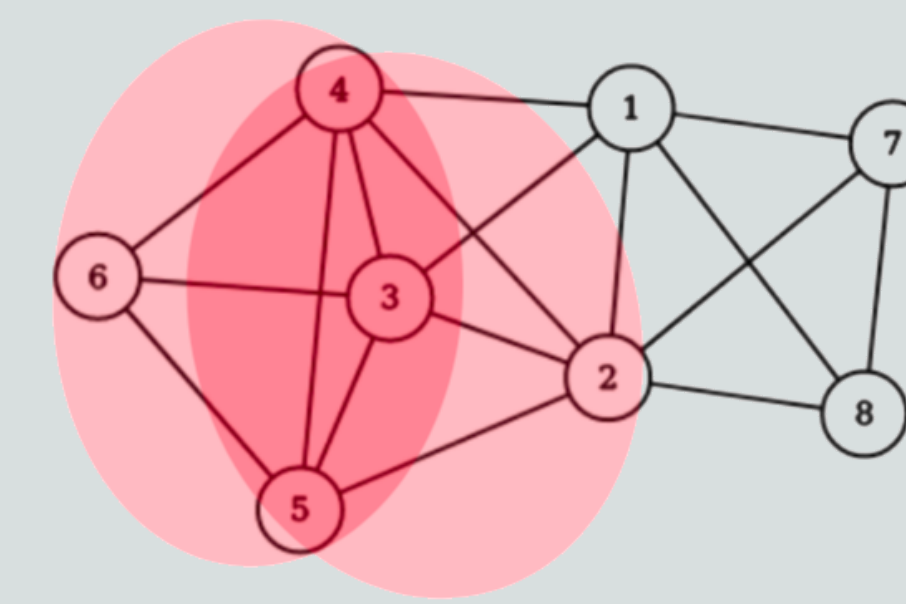
Palla et al., 2005

INTEREST

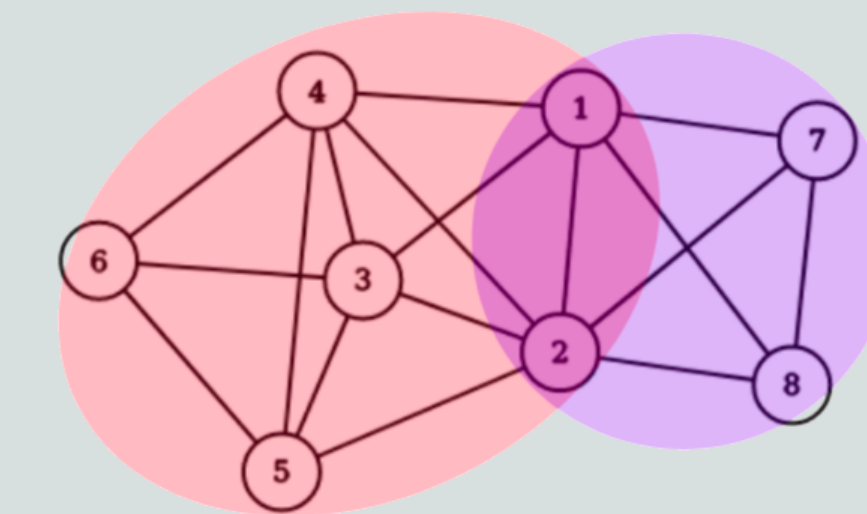
- Massive graphs : zoom in and out;
- Biological interactions;
- Content recommendation;
- ...

K-CLIQUE COMMUNITY

- **k-clique**: set of k not fully connected to each other.
- Two k-clique are **adjacent** if they share k-1 nodes.
- A **k-clique community** (CPM community), is the set of nodes of a maximal set of adjacent k-cliques.



Two adjacent k-cliques



k-clique communities

RELATED WORK

- Palla et al., 2005: first definition of CPM communities;
- Kumpula et al, 2008: solution based on k-clique enumeration;
- Reid et al, 2012: solution based on maximal clique enumeration;
- Gregori et al, 2013: parallel computation, based on maximal clique enumeration.

ALGORITHMS

METHOD

- Compute a stream of k-cliques in parallel;
- For each k-clique:
 - Find communities of each of its (k-1)-clique;
 - Merge them;
 - Add all new (k-1)-clique.

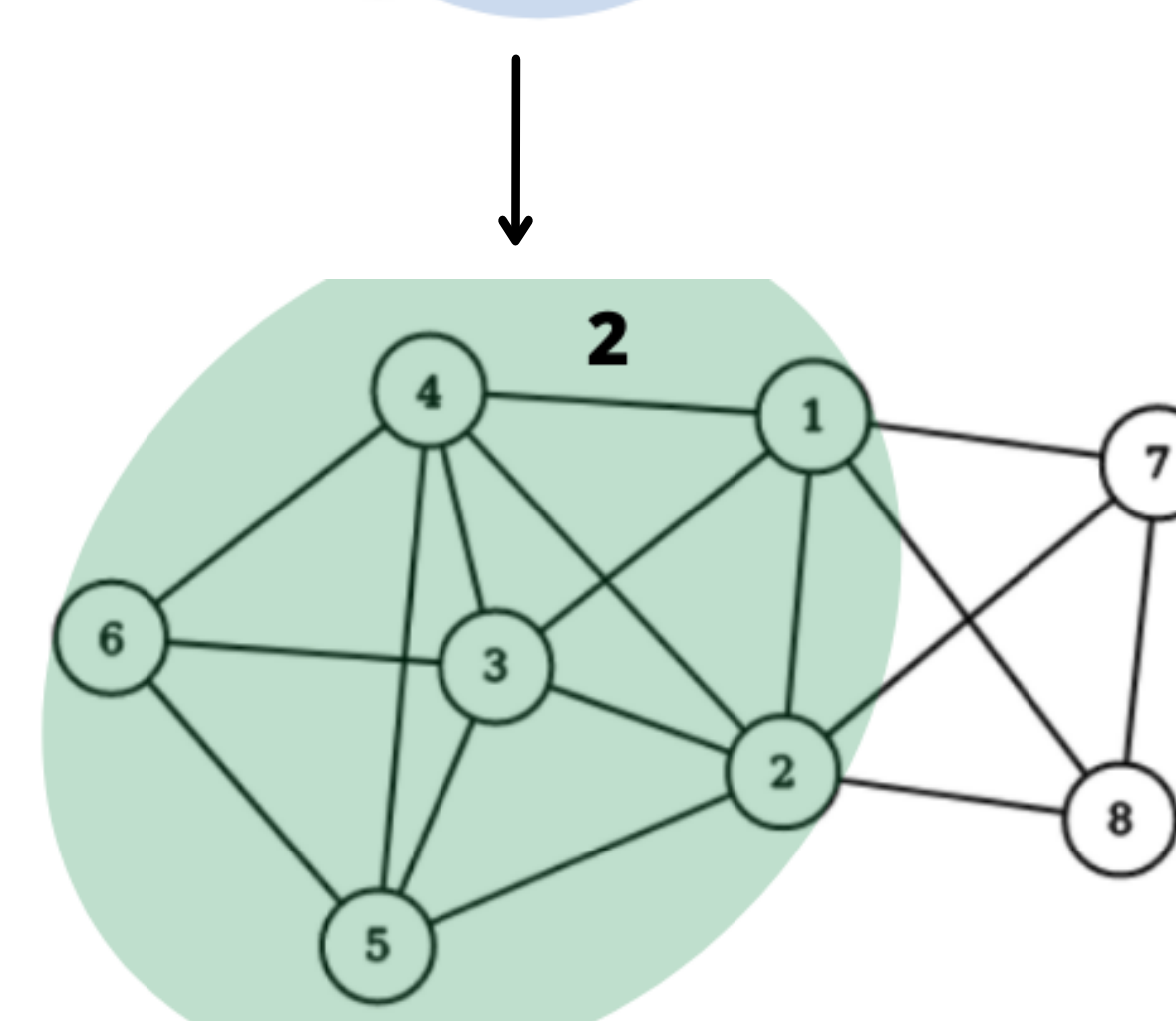
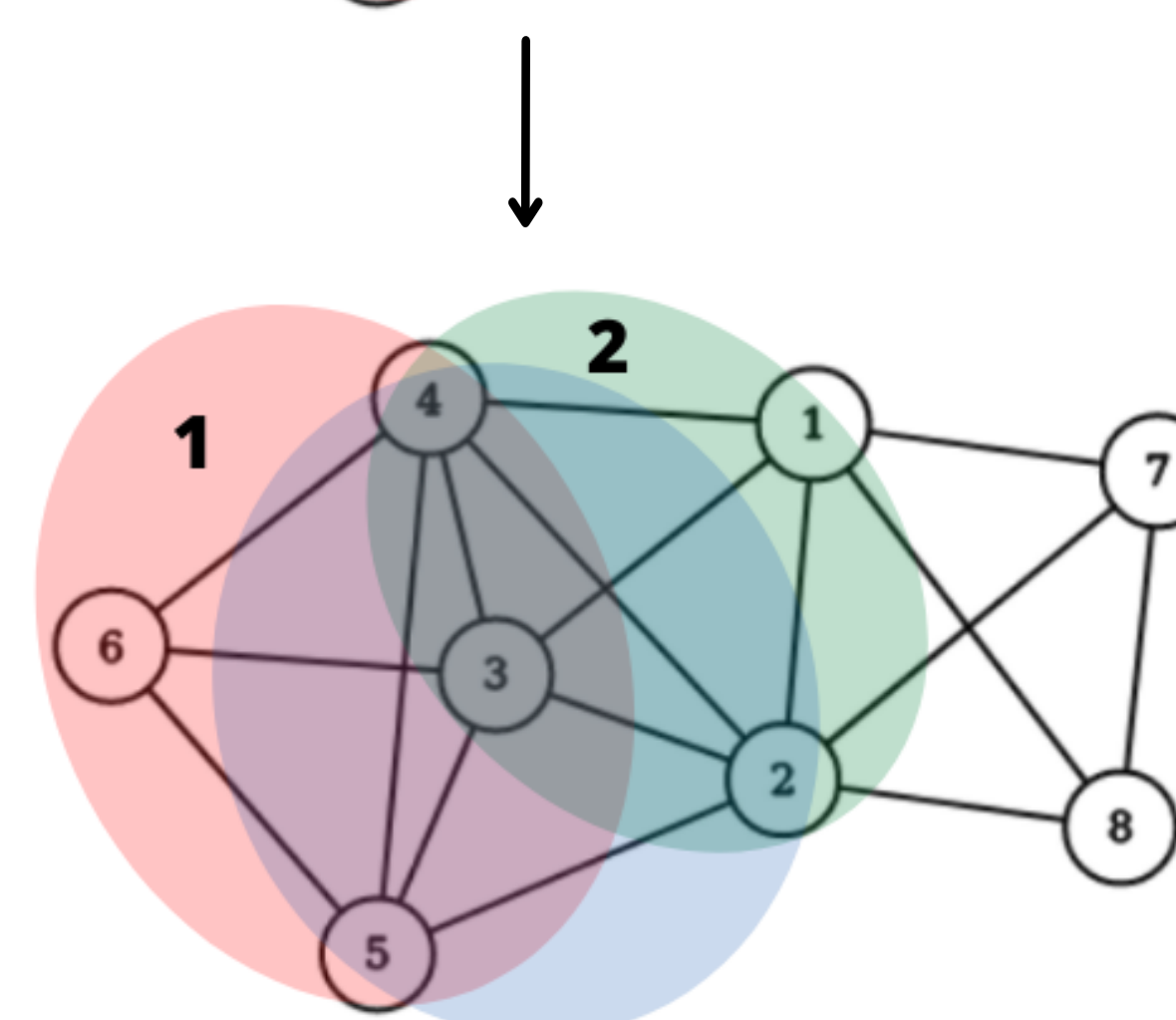
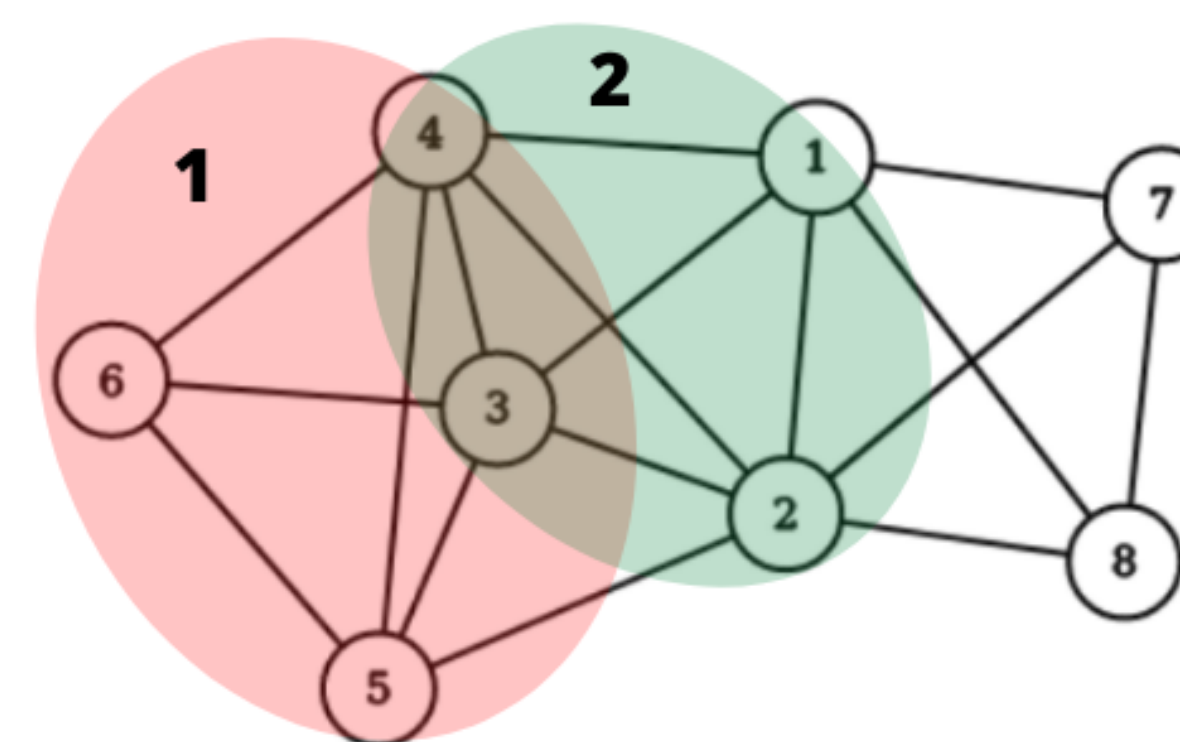
EXACT ALGORITHM

- Store a dictionary on (k-1)-cliques
- On the example:
 - New k-clique of the stream : (2,3,4,5)
 - Find:
 - (2,3,4) → 2
 - (2,3,5) → X
 - (2,4,5) → X
 - (3,4,5) → 1
 - Merge (Union-Find):
 - 2 ≅ 1
 - Add:
 - (2,3,5) → 2
 - (2,4,5) → 2

COMPLEXITY

Operations per k-cliques : $C_k \approx \mathcal{O}(k)$
(Find each (k-1)-clique)

EXAMPLE WITH K=4



MEMORY ISSUE OF EXACT ALGORITHM

Massive graphs : the larger k is, the more k-cliques there are.

HOW TO REDUCE MEMORY

- Exact algorithm : store all (k-1)-cliques;
- Approximate algorithm : store all z-cliques, for z < k-1.

APPROXIMATE ALGORITHM

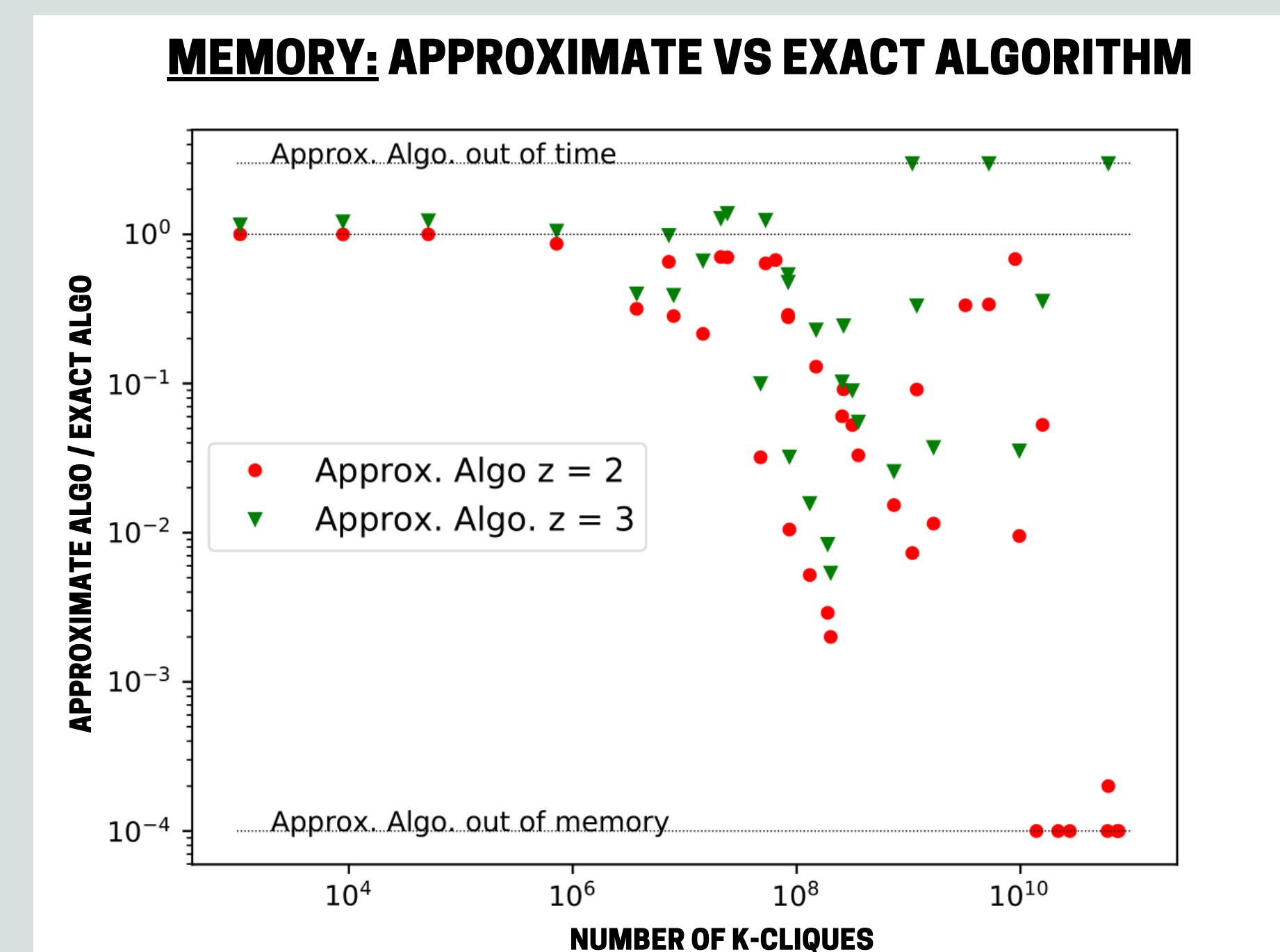
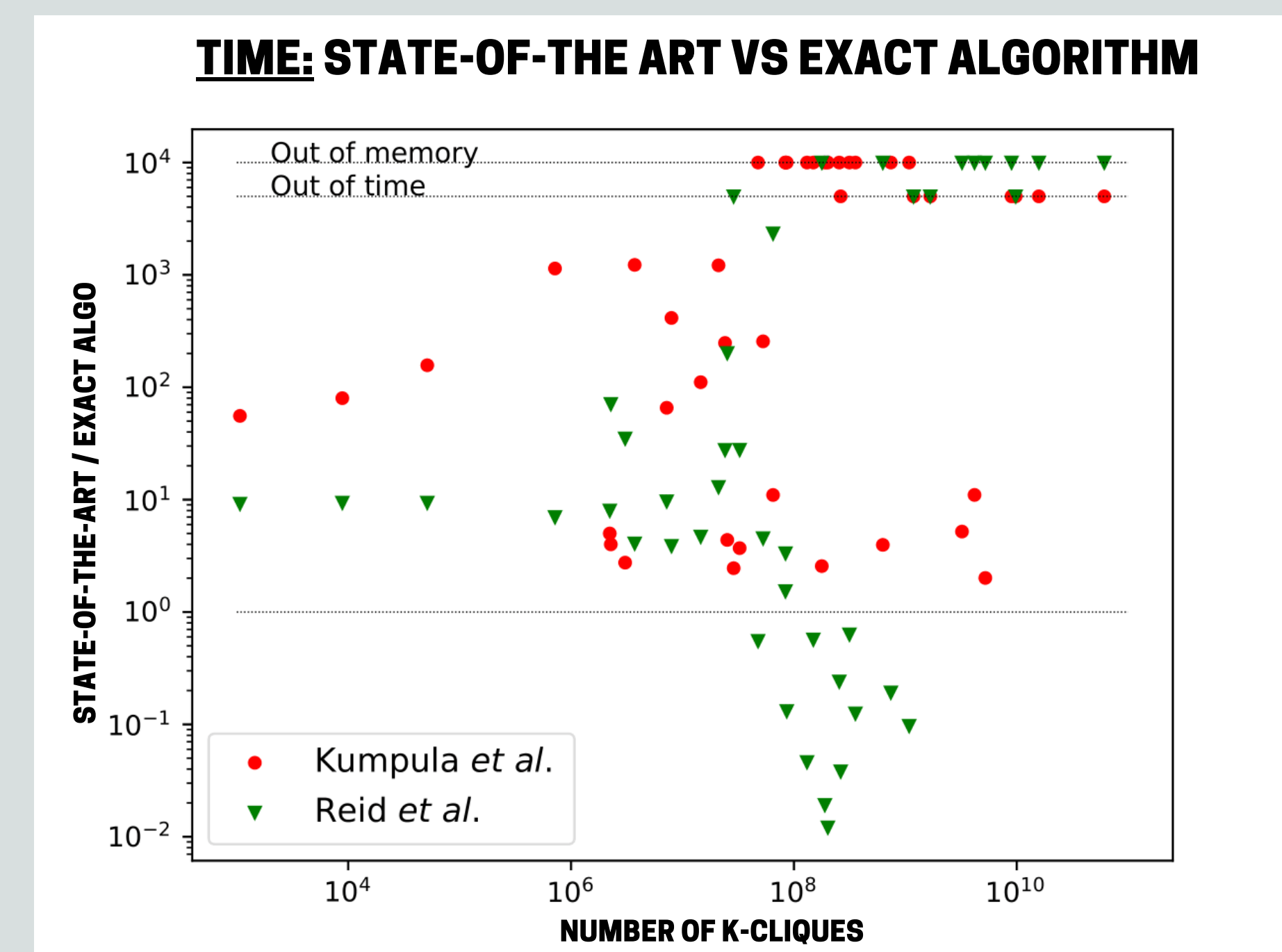
- Store a dictionary on z-cliques with z < k-1
- Difference with exact algorithm = Find
- On the example, with z=2 (store edges):
 - New k-clique of the stream : (2,3,4,5)
 - Find:
 - (2,3,4): 2
 - (2,3) → {2}
 - (2,4) → {2}
 - (3,4) → {1,2}
 - (3,4,5): 1
 - (3,4) → {1,2}
 - (3,5) → {1}
 - (4,5) → {1}

COMPLEXITY

Operations per k-cliques: $C_k \approx \mathcal{O}\left(k \cdot \binom{k-1}{z}\right)$
(Find each z-clique of each (k-1)-clique)

RESULTS

Communities of approximate algorithm are communities of exact algorithm, with some of them merged.



GOOD APPROXIMATION

Comparison Approximate VS Exact communities:

On all experiments:

- z=2: >93.8% of similarity, with mean=98.6%
- z=3: >99.5% of similarity, with mean=99.95%

Community comparisons with ONMI (MacAid et al. 2013)

CONCLUSION

- Exact algorithm: gain of time;
- Approximate algorithm: gain of memory.
 - best scale;
 - accurate approximate communities.